



**Description**

**[0001]** This invention relates to a packet-based radio telecommunications system, such as GSM (Global System for Mobile Communications), having improved traffic control, especially over the air interface, and relates especially to a Packet Control Unit (PCU) in the Base Transceiver Station (BTS) having improved interworking with Mobile Stations (MSs) and with the Serving GPRS Support Node (SGSN).

**[0002]** In a known packet-based system, a practical problem is that the PCU can only react on demand to packets received, and is not aware of traffic it will receive in future. The question is how to allocate resources on the air interface which carries GSM voice calls, control links, and timeslots. In current systems, the bit rate, the priority and the delay applied to each GPRS connection can be varied, but this requires considerable computing and messaging resources between the MS and the BTS, between the BTS and the PCU, and between the PCU and the SGSN.

**[0003]** It is an object of the invention to optimise the usage of radio resources without increasing the computing and signalling loads on the network elements or the MS.

**[0004]** A further object is to assist in maintaining the blocking rate of the system at an acceptable level. The blocking rate is the proportion of users who are trying to access the system but cannot immediately be served; the rate can never be zero, which implies infinite resources, and practical experience shows that a blocking rate of 5% generates a substantial level of caller complaints. A blocking rate of 2% has been found to be acceptable, and a maximum number of GPRS transmissions can then be served.

**[0005]** According to the invention, in a radio telecommunications system comprising a node operating the leaky bucket algorithm to control flow of packet data to a local packet control unit controlling at least one radio telecommunications cell, the algorithm having a variable maximum bucket size Bmax and a variable leak rate R, and there being a virtual queue of data formed in the local control unit, a method of controlling data flow from the node to the local control unit characterised by :-

providing lower and upper trigger lengths on the virtual queue;  
sensing when the virtual queue reaches the upper trigger; and  
sending to the node a request to alter the value of R so as to reduce the rate of data flow to the local control unit.

**[0006]** Also according to the invention, a packet control unit for a radio telecommunications system comprises a down link requests scheduler; a base station system virtual connection flow controller; a radio link control /medium access control blocks scheduler; a mobile system flow controller; a first data transmission means to exchange data with a plurality of mobile systems; and a second data transmission means to exchange data with the serving support node of the radio telecommunications system; the base station system virtual connection flow controller having a data signalling queuing means provided with upper and lower timeslot triggers, characterized in that said data signalling queuing means is further provided with upper and lower BVC flow control triggers, both at a greater queue length than the timeslot triggers.

**[0007]** Preferably a data rate limitation request message is sent to the serving support node only when the queue length is between the additional upper and lower triggers, such a condition being referred to as a congestion condition.

**[0008]** Optionally, in a congestion condition, the queue length is arranged to oscillate about a position midway between the upper and lower additional triggers.

**[0009]** The invention will now be described by way of example with reference to the accompanying drawings in which:-

**[0010]** Figure 1 is a schematic illustration of the base station part of the GPRS (General Packet Radio Service) radio telecommunications system; and Figure 2 is a schematic illustration of the packet control unit of a radio telecommunications system.

**[0011]** In Figure 1 a GPRS system 10 comprises a GPRS Backbone System (GBS) 12 containing a Serving GPRS Node (SGSN) 14 which is connected by an interface link Gb 16 to a Packet Control Unit (PCU) 18 within a Base Station System (BSS) 20. The PCU 18 is connected to a number of Mobile Systems (MS) 22 through an interface link Um 24.

**[0012]** In Figure 2 the PCU 18, the link Gb 16, and the link Um 24 are shown. Between the links 16 and 24 there is a Down Link (DL) requests scheduler 26, a MS flow control unit 28, a Base Station System Virtual Connection (BVC) queue scheduler 30; a BVC flow controller 32, and a Radio Link Control/Medium Access Control (RLC/MAC) blocks scheduler 34.

**[0013]** Although the uplink is not illustrated as it is not actively involved in application of the present invention, it will naturally exist.

**[0014]** In normal operation at normal traffic flow rates, a new call to a MS 22 is received over the Gb link 16, and the first Protocol Data Unit DL-UNITDATA PDU passes to the queue 38 of the DL requests scheduler 26. The scheduler 26 instructs the RLC/MAC blocks scheduler 34 to allocate capacity, and a Temporary Block Flow (TBF) queue, such as queue 42, is set up for the called mobile. Signals pass over the Um link 24 to the called mobile, which returns acknowledgement signals, ACK, over link 24 to the queue 42. The scheduler 34 fetches the next DL-UNITDATA PDU

which, because the MS is now known, passes through the BVC queue scheduler 30 into the BVC queue 36, and then to the appropriate TBF queue, queue 42 in this example.

**[0015]** The UNITDATA PDU corresponds to many messages and acknowledgements, and use of the RLC/MAC scheduler reduces the volume for the MS.

**[0016]** The scheduler 26 assigns logical resources to one of the TBF queues 40, 42, and instructs the RLC/MAC block scheduler, message(2) - (2). The processor 28 reads from the second DL-UNITDATA PDU the multi slot capacity of the mobile.

**[0017]** One TBF queue is allocated to each active MS. Every MS with an active TBF queue has two logical queues; one is for Logical Link Layer (LLC) data PDUs, and one is for LLC signaling PDUs. The length of the BVC queue 36 is the sum of all the MS queues including LLC data and signaling.

**[0018]** The BVC queue 36 is provided towards its output end with two timeslot triggers TS(l) and TS(u) which function in the conventional way, that is, if the queue 36 exceeds the upper timeslot trigger TS(u), the DL requests scheduler 36 sends a request via interface Gb 16 to the SGSN 14 for an additional timeslot to be allocated. If an additional timeslot is allocated, the queue shortens as traffic throughput is increased, and the queue length falls below the upper trigger TS(u). The signal from the queue 36 to the scheduler 26 is indicated as message (3) - (3).

**[0019]** In the arrangement of Figure 2, the major components illustrated perform their conventional functions as follows:-

**[0020]** The DL requests scheduler 26 allocates resources, ie. it decides which is the next MS 22 to be connected over the link 24; it requests additional timeslots when the timeslot trigger TS(u) is activated on connection (3) - (3), and it returns timeslots when they are no longer required.

**[0021]** The RLC/MAC blocks scheduler 34 allocates capacity, i.e. it allocates timeslots in response to instructions from the DL scheduler 36, message (2) - (2); it fetches the next DL PDU when the TBF queue 40, 42 for that MS is below a predetermined threshold (precedence is given to LLC signaling PDUs). The scheduler 34 also divides the LLC PDU into blocks, sets up transmission windows and retransmissions of blocks for each TBF queue 40, 42, and drops the TBF after a predetermined number of resent transmissions; on termination of a call it also signals the end of a TBF queue to the DL scheduler 36, message (1) - (1).

**[0022]** The DL requests scheduler 26 and the RLC/MAC blocks scheduler 34 are related as follows; the scheduler 26 assigns logical resources to each TBF queue 40, 42 on a per timeslot basis, ie. an allocation table is set up, while the scheduler 34 dynamically maps these logical resources to physical resources. Looked at another way, the scheduler 26 decides which mobile 22 will next be connected, and the scheduler 34 decides which block of information is sent to which mobile.

**[0023]** The BVC queue scheduler 30 selects the appropriate queue for a mobile system. It directs each DL-UNITDATA PDU to that queue. Each MS queue is divided into an MS data queue and an MS signaling queue - conveniently via use of two pointers in a common buffer.

**[0024]** Data flow over the Gb link 16 is controlled by use of the leaky bucket algorithm run within the SGSN 14 (Figure 1). Values of Bmax, the maximum bucket capacity and R the leak rate, are calculated by both the BVC flow control processor 32 and the MS flow control processor 28 in different circumstances, and the values sent over the Gb link 16.

**[0025]** The following assumptions are made about the SGSN 14;

1 The MS flow control values sent to the SGSN by the MS flow control processor 28 are valid until the SGSN 14 receives a new valid MS flow control message or until the hysteresis timer Th expires. The PCU 18 knows the timer value which the SGSN is using, and is arranged to send a new MS flow control message every Th seconds less a short tolerance time.

2 If the SGSN 14 has not received an MS flow control message for a particular MS, the SGSN uses the default flow control values as sent in every BVC flow control message. The SGSN never uses internally generated MS flow control values, when the initial values are sent in the BVC flow control message.

3 The SGSN 14 sends the MS Radio Access Capability (a class mark) in each DL-UNITDATA PDU.

4 The SGSN 14 keeps (logical) separated queues for both MS flow control and BVC flow control.

**[0026]** Turning now to detailed consideration of the inventive feature, the BVC queue 36 is provided with two additional triggers, BVC(l) and BVC(u), at greater queue lengths than the timeslot triggers TS(l) and TS(u). Between BVC(l) and BVC(u) there is an intermediate queue length which is referenced Bmax in Figure 2.

**[0027]** Suppose the upper timeslot trigger TS(u) is reached, and an additional timeslot is requested, but is not received; this may occur if all timeslots are occupied by voice calls, or if the PCU has already reached the maximum number of timeslots available, i.e., a whole cell is occupied so no additional timeslot can be allocated.

**[0028]** In these conditions the queue 36 is likely to grow and to exceed the position of BVC(l) trigger. When this trigger is reached, conditions will be referred to as congestion conditions.

**[0029]** If the queue continues to lengthen and reaches the upper trigger BVC(u), the BVC flow controller 32 then

calculates the desired bit rate to be received from the SGSN using as inputs the queue length and the former bit rate value sent in the BVC flow control message, and sends a BVC flow control message over connection 46 through the interface Gb to the SGSN, requesting that data is sent at a lower bit rate. When the SGSN responds, the queue 36 shortens and may fall below the BVC lower trigger, when a further message is sent, increasing the bit rate.

**[0030]** Subsequently, either an additional timeslot may be allocated to the PCU, or traffic eases; in either case the length of the queue 36 will decrease and may go out of congestion conditions by falling below the lower trigger BVC (1), then below the trigger TS(u). The timeslot triggers then function in the conventional way.

**[0031]** It will be clear that the BVC triggers operate in real time, and that dynamic reaction to traffic conditions is achieved.

**[0032]** An advantage of this use of two sets of upper and lower triggers is that control of the SGSN 14 by the PCU 18 is light; the SGSN has priority of control in most circumstances, e.g. on GPRS delay on the down link, because it has information not available to the PCU. Further, the BVC triggers do not interfere with the timeslot triggers, by reason of the small gap between TS(u) and BVC(l).

**[0033]** Although the request queue 38, the BVC queue 36 and the TBF queues 40, 42 are shown as logically separate queues, in practice all three queues share the same physical space, and can be implemented as separate pointers in that space.

**[0034]** Considering now the operation of the four triggers on the BVC queue in more detail; the general concept is that of capacity on demand.

**[0035]** The DL scheduler 26 requests a new timeslot when either:

- a) a certain ratio of served TBFs for a timeslot is reached, e.g. when there are already five active TBFs for each timeslot and a new request is received; or
- b) the BVC queue reaches the upper timeslot trigger TS(u), or
- c) when both conditions apply.

**[0036]** The objective of condition b), i.e. the position at which the upper timeslot trigger is set, is to have in the queue a new LLC-PDU for each TBF before the previous one has been transmitted; for example, to have an average of 1.5 LLC-PDUs for each TBF, which means that one LLC-PDU is queued and half of the previous one has been transmitted.

**[0037]** A typical value of the upper timeslot trigger is TS(u) \* number of timeslots. The TS(u) value is chosen in order to limit the delay in the BSS avoiding at the same time the risk of early TBF termination.

**[0038]** As referred to above, the BVC flow control processor 32 sends intermittent flow control messages through Gb link 16 to the SGSN 14, the messages being the maximum BVC bucket size BVCBmax and the BVC leak rate R for use in the leaky bucket algorithm run by the SGSN 14.

**[0039]** In the arrangement according to the invention, the upper timeslot trigger TS(u) is activated before BVC control applies. It is now suggested that BVC Bmax calculated by the controller 32 is made proportional to the number of allocated timeslots, i.e., it is unchanged so long as the number of timeslots/packet data channels is unchanged. It is also now possible for real-time BVC flow control to be exerted only by changing the value of R.

**[0040]** An objective of the BVC flow control algorithm is to have an amount of data equal to Bmax in the BVC queue 36 during congestion conditions; effectively, the queue length oscillates about the length equal to Bmax. As a result of this, the TS(u) remains active and the system keeps on asking for a further slot or further capacity, and the queue length is controlled, and the delay in the BSS is also controlled.

**[0041]** Considering now some appropriate values; the distance between TS(u) and BVC(1) is set equal to the average size of an LLC PDU, and so is the distance between the TS(l) and BVCBmax.

**[0042]** A value of BVCBmax can then be set as :

$$\text{BVC Bmax (for a PDCH)} = (\text{TS(u)} + \text{average LLC PDULength} * 2) + 10\%$$

(10% is added to maintain reasonable tolerance)

**[0043]** As stated above, an objective of the BVC flow control algorithm run by the BVC flow control processor 32 is to have the amount of data equal to Bmax in the BVC queue 36 during congestion conditions. The results are that:

- 1) Controlling the queue so that it is no longer than Bmax (with some tolerance) means that the delay introduced in the BSS on average is no longer than Bmax/((average bit rate for a timeslot)\* number of timeslots). Bmax is proportional to the number of timeslots, so that delay = Bmax for a timeslot/average bit rate for a timeslot
- 2) Controlling the queue so that it is no shorter than Bmax (with some tolerance) means that the timeslot upper trigger for requesting a new timeslot is always active.

[0044] Since the value of Bmax is proportional to the number of timeslots, the BVC flow control can be exerted by means of the leak rate R.

[0045] The maximum value of R, i.e., the maximum bit rate which a BVC can support is equivalent to the maximum capacity of a PDCH. The minimum R, i.e., the minimum bit rate which a BVC can support is equivalent to the minimum capacity of a PDCH. The optimal value of R will be 10 or 20% lower than the maximum.

[0046] The actual value of R will vary between the maximum and minimum, and it is suggested that for every change in the number of channels, the value of R sent in the BVC flow control message is the optimal R multiplied by the number of channels, thus:

$$\text{BVC R} = \text{optimal value of R} * (\text{number of PDCHs})$$

[0047] If the real bit rate is lower, the PCU 18 will provide a new flow control message. Using the value set out above, the PCU can be sure that the SGSN 14 will send enough data to activate the upper timeslot trigger.

[0048] Turning now to the positions of the BVC upper and lower triggers, these can be set as the value of Bmax plus and minus 20%.

[0049] When the trigger is triggered, the PCU 14 must react by sending a value of R which allows the queue 36 to equal the size of Bmax within x seconds, where x is called the converging period; the converging period should never be shorter than the C timer, but should not be too short in order to avoid a sharp movement in the opposite direction; also the period should not be too long.

[0050] The BVC flow control process uses the old value of R and the elongation in the queue length (queue length minus Bmax) as a basis for calculating the new value of R; this has the advantage of avoiding any reference to the actual bit rate over the Um interface 24. The new value of R should allow the queue to return to equal Bmax within the converging period, but must never be higher than a maximum R or lower than the minimum R in order to avoid sharp changes in queue length. A new value for R is given by:

$$\text{new R} = \max[\min[\text{old R} + ((\text{Bmax} - B)^* 8 / (\text{converging period})), \text{Maximum R}], \text{Minimum R}]$$

where  $(\text{Bmax} - B)^* 8$  is the queue elongation in bits.

[0051] Consider now a change in the number of PDCHs. If a channel is added, the PCU 14 will have more capacity on the link 24, so there will be no problems in managing the queues. If a channel is released, the PCU has ensured both that the queue is really short (below the lower timeslot trigger) and that the arrival rate from the SGSN 14 can be served with one channel fewer. In both cases, the trigger for a new BVC flow control message must not be too fast so that the situation stabilizes. Waiting for some seconds and then reevaluating B and R reduces the risk of starting an elongation in queue length.

[0052] Further, it is advantageous to avoid sending a new BVC flow control message too quickly after a previous message. A BVC flow control check timer can be used, equal to or higher than the C timer, and preferably at least half of the converging period. The arrangement is such that while this timer is running, no check is performed on BVC triggers.

[0053] Taking an overall view, the BVC flow controller 32 operates in two phases:

A) Initialization: this is necessary when the BVC connection is created, after a reset or an unblocking procedure. The PCU 18 sends a BVC flow control message the SGSN 14; at this point the PCU 18 has no timeslots, thus flow control values are very close to zero.

B) Real time: the PCU needs to control the actual bit rate from the SGSN 14 on the Gb link 16, and the bit rate is adjusted by sending new values of B and R to the SGSN. Such a message is sent when either:

- a) the PCU 18 adds or releases timeslots, when the BVC flow control is used to avoid buffer overload or lack of data to send. The PCU can use fixed values related to the number of allocated timeslots, so no heavy computation is required OR
- b) congestion conditions apply, i.e., the BVC upper trigger has been activated as described above.

[0054] The DL scheduler 26 releases a timeslot when it is no longer needed. For example, if the PCU has a single timeslot, the last timeslot in the cell can be released after the last TBF has been deactivated, plus a hysteresis time, provided there are no more LLC-PDUs waiting to be transmitted. Consideration also needs to be given to the distribution of active TBFs; it may be more efficient to retain the timeslot and reallocate TBFs to timeslots to give a more even

distribution, i.e., avoiding eight TBFs in one timeslot and one TBF in another timeslot.

[0055] By application of the invention, the operation of the SGSN 14 is limited only when absolutely necessary, by means of dynamic resource allocation. Only whole timeslots are used, never part of a timeslot. Such a timeslot is allocated dynamically in both the GSM voice services and in GPRS, and by application of the invention, the allocation is made efficient.

## Claims

1. In a radio telecommunications system comprising a node (14) operating the leaky bucket algorithm to control flow of packet data to a local packet control unit (18) controlling at least one radio telecommunications cell, the algorithm having a variable maximum bucket size Bmax and a variable leak rate R, and there being a virtual queue (36) of data formed in the control unit (18), a method of controlling data flow from the node (14) to the control unit **characterised by :-**

providing lower and upper triggers on the virtual queue (36);  
sensing when the virtual queue reaches the upper trigger; and  
sending to the node(14) a request to alter the value of R so as to reduce the rate of data flow to the local control unit.

2. A method according to Claim 1 further comprising:-

sensing when the virtual queue (36) reaches the lower trigger; and  
sending to the node (14) a request to alter the value of R such that the length of the virtual queue is maintained close to the value of Bmax.

3. A method according to Claim 1 or Claim 2 in which the value of Bmax is proportional to the number of timeslots allocated by the node (14) to the local control unit (18).

4. A method according to any one of Claims 1, 2 or 3 further comprising the step of requesting said node 14 to allocate an additional time slot to the local control unit when the virtual queue is longer than the upper trigger and/or the number of temporary block flow queues per time slot allocated to the local control unit exceeds a predetermined threshold.

5. A method according to any one of Claims 1, 2, 3 and 4 comprising the further steps of the local control unit sensing when the last temporary block flow queue for a timeslot has been released and whether that timeslot is the last GPRS timeslot in the cell, and the local control unit then releasing said last timeslot.

6. A method according to claim 5 in which when there is more than one GPRS timeslot allocated to said local control unit, the local control unit releases a timeslot when it senses that:-

the queue length is lower than the timeslot lower trigger,  
and the number of temporary block flow queues per timeslot, when calculated on the basis of one timeslot less than the number currently allocated, is lower than a predetermined threshold,  
and the arrival rate of data from the node is lower than the arrival rate which can be supported by one timeslot less than the number currently allocated.

7. A method according to any one of claims 1 to 6 in which the local control unit is further provided with upper and lower timeslot triggers at a queue length lower than said upper and lower BVC triggers, and when the queue length is greater than the upper timeslot trigger, and a new timeslot cannot be allocated, the local control unit maintains the length of the virtual queue at a value higher than the upper timeslot trigger.

8. A packet control unit (18) for a radio telecommunications system (10) comprises a down link requests scheduler (26); a base station system virtual connection flow controller (32); a radio link control/medium access control blocks scheduler (34); a mobile system flow controller (28); first data transmission means (24) to exchange data with a plurality of mobile systems (22); and second transmission means (16) to exchange data with a serving support node (14) of the radio telecommunications system; the base station system virtual connection flow controller (32) having a data signalling queuing means provided with upper and lower timeslot triggers, TS(u), TS(l) characterized

in the said data signalling queuing means is further provided with upper and lower BVC flow control triggers BVC (u), BVC (l), both at a greater queue length than the timeslot triggers.

5     **9.** A packet control unit (18) according to Claim 8 in which the base station system virtual connection flow controller (32) is arranged to send a data limitation request message to a serving node (14) of the system only when the packet control unit is in congestion conditions as hereinbefore defined or when there is a change in the number of allocated timeslots.

10    **10.** A packet control unit (18) according to claim 8 or claim 9 in which when the queue length is above the lower additional trigger BVC(l), the queue length is arranged to oscillate about a value Bmax between the upper and lower triggers.

15

20

25

30

35

40

45

50

55

FIG. 1

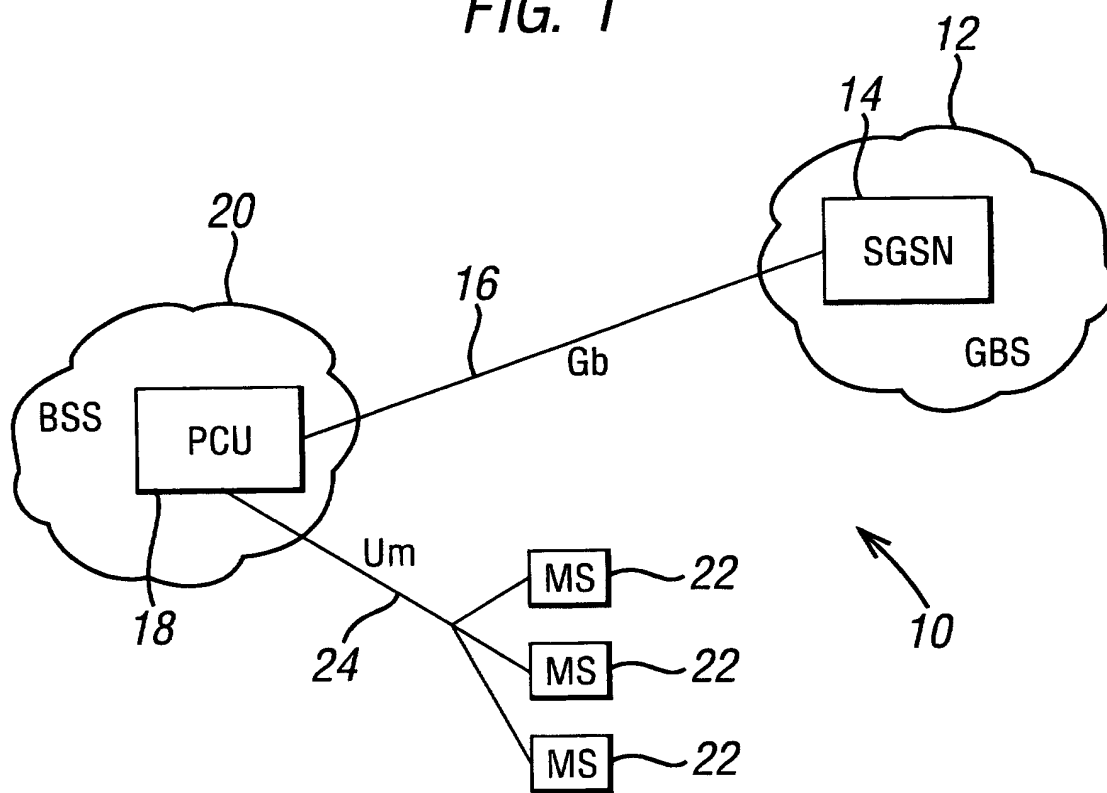
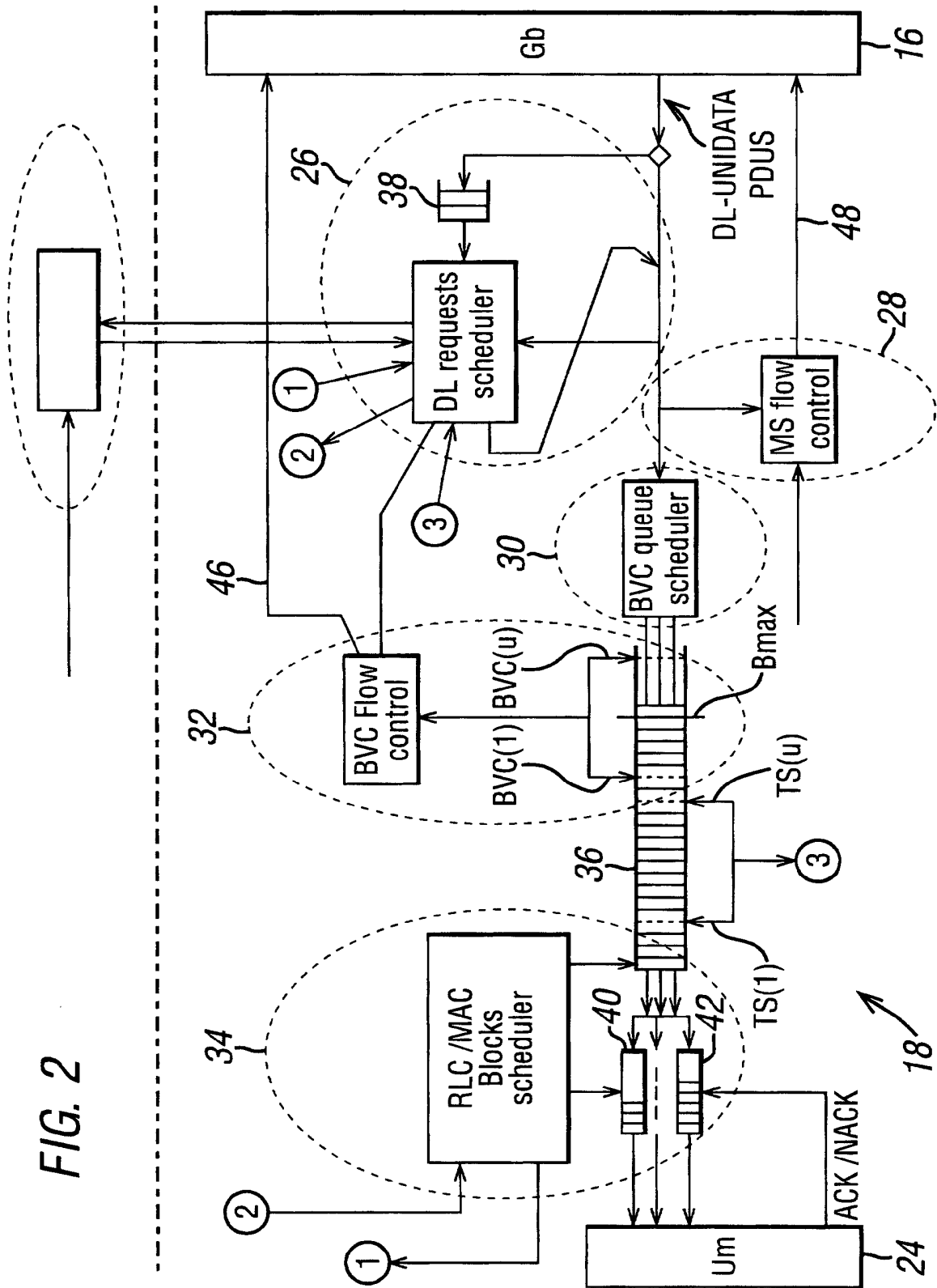




FIG. 2





European Patent  
Office

## EUROPEAN SEARCH REPORT

Application Number  
EP 00 30 1874

| DOCUMENTS CONSIDERED TO BE RELEVANT   |  |  |  |
|---|--|--|--|
| Category  | Citation of document with indication, where appropriate, of relevant passages  | Relevant to claim  | CLASSIFICATION OF THE APPLICATION (Int.Cl.7) |
| A   | US 5 696 764 A (SOUMIYA ET AL.)<br>9 December 1997 (1997-12-09)<br>* abstract; figures 1A-2,11,14,15 *<br>* column 14, line 27 - column 15, line 11 *<br>* column 26, line 47 - line 50 *<br>--- | 1,2,8  | H04Q7/22                                     |
| A   | US 5 802 310 A (RAJARAMAN)<br>1 September 1998 (1998-09-01)<br>* abstract; figures 5-7 *<br>* column 4, line 9 - line 50 *<br>---  | 1,2,8  |  |
| A   | WO 97 39555 A (DSC COMMUNICATIONS A/S)<br>23 October 1997 (1997-10-23)<br>* page 5, line 20 - page 6, line 5; claims 1,7 *<br>---  | 1,2,8  |  |
| A   | EP 0 874 532 A (MMC NETWORKS, INC.)<br>28 October 1998 (1998-10-28)<br>* figures 8,9 *<br>* page 4, line 26 - line 36 *<br>---   | 1,2,8  |  |
| A   | US 5 914 936 A (HATONO ET AL.)<br>22 June 1999 (1999-06-22)<br>* abstract; claims 1-4; figure 4 *<br>---   | 1  | H04Q<br>H04L                                 |
| A   | US 5 570 360 A (KLAUSMEIER ET AL.)<br>29 October 1996 (1996-10-29)<br>* abstract; figures 7-10A,11,12 *<br>* column 8, line 59 - column 9, line 8 *<br>---                                       | 1  |  |
| A   | EP 0 647 081 A (K.K.TOSHIBA)<br>5 April 1995 (1995-04-05)<br>* abstract; figures 4,5 *<br>* column 11, line 33 - line 55 *<br>-----  | 1  |  |
| The present search report has been drawn up for all claims  |  |  |  |
| Place of search<br><b>BERLIN</b>  |  | Date of completion of the search<br><b>3 August 2000</b>   | Examiner<br><b>Danielidis, S</b>             |
| <b>CATEGORY OF CITED DOCUMENTS</b><br>X : particularly relevant if taken alone<br>Y : particularly relevant if combined with another document of the same category<br>A : technological background<br>O : non-written disclosure<br>P : intermediate document |  | T : theory or principle underlying the invention<br>E : earlier patent document, but published on, or after the filing date<br>D : document cited in the application<br>L : document cited for other reasons<br>& : member of the same patent family, corresponding document |  |

EPO FORM 1503 03/82 (P04C01)

**ANNEX TO THE EUROPEAN SEARCH REPORT  
ON EUROPEAN PATENT APPLICATION NO.**

EP 00 30 1874

This annex lists the patent family members relating to the patent documents cited in the above-mentioned European search report.  
The members are as contained in the European Patent Office EDP file on  
The European Patent Office is in no way liable for these particulars which are merely given for the purpose of information.

03-08-2000

| Patent document<br>cited in search report | Publication<br>date | Patent family<br>member(s) | Publication<br>date |
|---|---------------------|----------------------------|---------------------|
| US 5696764 A                              | 09-12-1997          | WO 9503657 A               | 02-02-1995          |
|   |                     | US 5936958 A               | 10-08-1999          |
| US 5802310 A                              | 01-09-1998          | NONE                       |                     |
| WO 9739555 A                              | 23-10-1997          | DK 42896 A                 | 13-10-1997          |
|   |                     | AU 2380897 A               | 07-11-1997          |
|   |                     | EP 0894384 A               | 03-02-1999          |
| EP 874532 A                               | 28-10-1998          | US 6014367 A               | 11-01-2000          |
|   |                     | JP 10327170 A              | 08-12-1998          |
| US 5914936 A                              | 22-06-1999          | JP 9307561 A               | 28-11-1997          |
| US 5570360 A                              | 29-10-1996          | AU 5183196 A               | 08-10-1996          |
|   |                     | CA 2211077 A               | 26-09-1996          |
|   |                     | EP 0815674 A               | 07-01-1998          |
|   |                     | WO 9629805 A               | 26-09-1996          |
| EP 647081 A                               | 05-04-1995          | JP 7123102 A               | 12-05-1995          |
|   |                     | KR 146020 B                | 17-08-1998          |
|   |                     | US 5694390 A               | 02-12-1997          |
|   |                     | US 5835484 A               | 10-11-1998          |

EPO FORM P0459

For more details about this annex : see Official Journal of the European Patent Office, No. 12/82